**Pune Institute of Computer Technology**

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**Department of Computer Engineering**

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# **“Different exact and approximation algorithms for Travelling-Sales-Person Problem”**

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**Savitribai Phule Pune University**

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**In**

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**By**

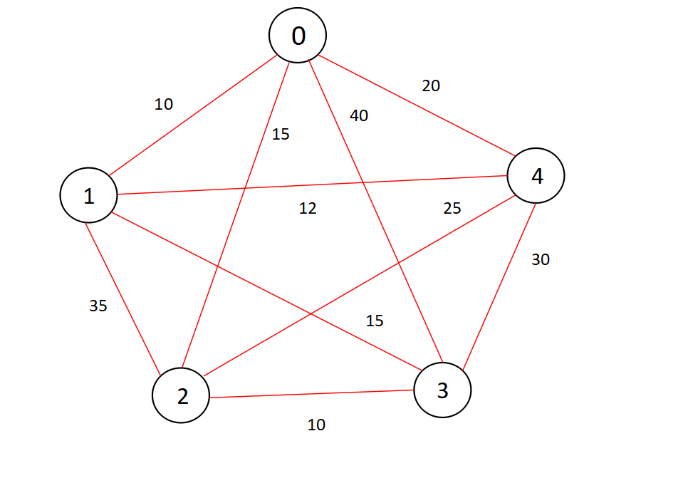
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**Problem Statement:**

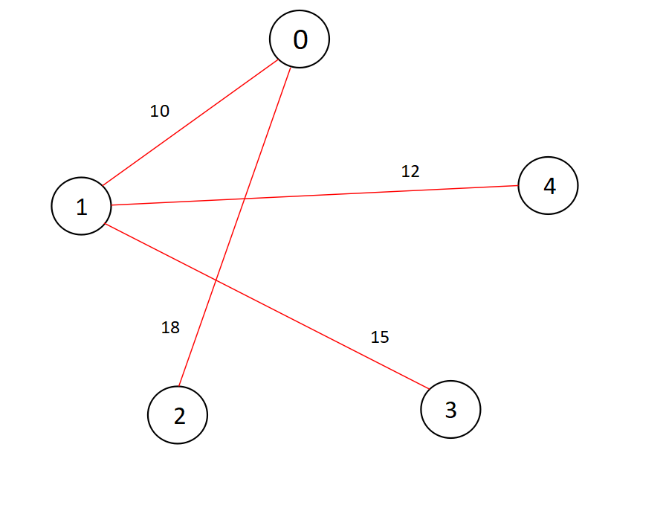
Different exact and approximation algorithms for Travelling-Sales-Person Problem

**Description:**

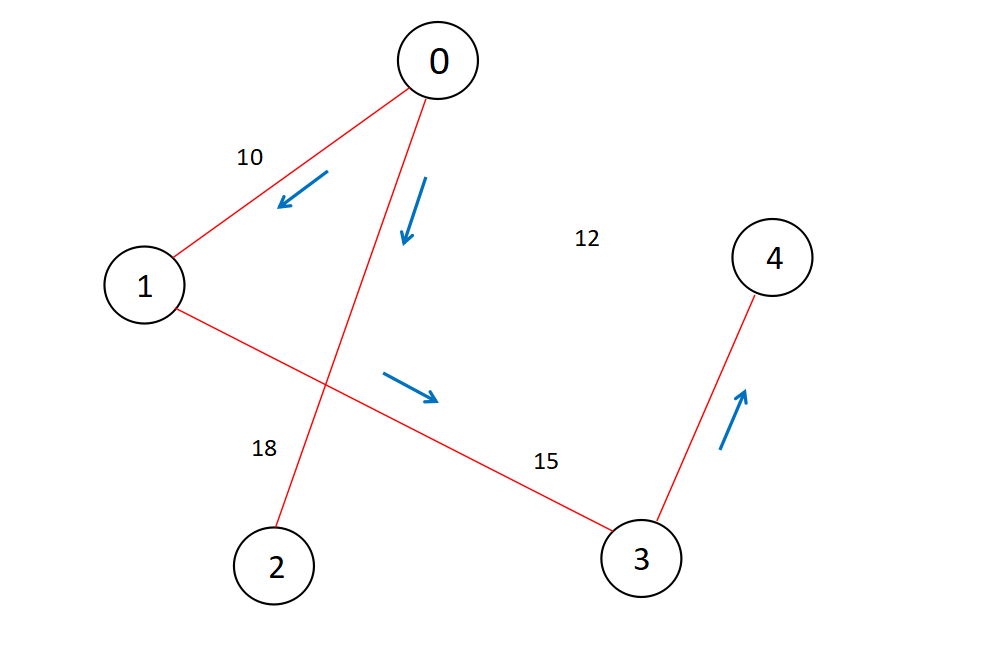
1. **Travelling Sales Person Problem**
   * Travelling Salesman Problem is based on a real life scenario, where a salesman from a company has to start from his own city and visit all the assigned cities exactly once and return to his home till the end of the day.
   * The exact problem statement goes like this, **"Given a set of cities and distance between every pair of cities, the problem is to find the shortest possible route that visits every city exactly once and returns to the starting point."**
   * There are two important things to be cleared about in this problem statement,
     + **Visit every city exactly once**
     + **Cover the shortest path**
2. **Designing the code:**
   * Step - 1 - Constructing The Minimum Spanning Tree
     + Creating a set mstSet that keeps track of vertices already included in MST.
     + Assigning a key value to all vertices in the input graph. Initialize all key values as INFINITE. Assign key value as 0 for the first vertex so that it is picked first.
     + [The Loop] While mstSet doesn’t include all vertices
     + Pick a vertex u which is not there in mstSet and has minimum key value.(minimum\_key())
     + Include u to mstSet.
     + Update key value of all adjacent vertices of u. To update the key values, iterate through all adjacent vertices. For every adjacent vertex v, if weight of edge u-v is less than the previous key value of v, update the key value as weight of u-v.
   * Step - 2 - Getting the preorder walk/ Defth first search walk:
     + Push the starting\_vertex to the final\_ans vector.
     + Checking up the visited node status for the same node.
     + Iterating over the adjacency matrix (depth finding) and adding all the child nodes to the final\_ans.
     + Calling recursion to repeat the same.
3. **Example:**
   * Let's have a look at the graph(adjacency matrix) given as input



* + After performing step-1, we will get a Minimum spanning tree as below



* + Performing DFS, we can get something like this



**Code:**

/\*

#include <bits/stdc++.h>

using namespace std;

// Number of vertices in the graph

#define V 5

// Dynamic array to store the final answer

vector<int> final\_ans;

int minimum\_key(int key[], bool mstSet[])

{

int min = INT\_MAX, min\_index;

for (int v = 0; v < V; v++)

if (mstSet[v] == false && key[v] < min)

min = key[v], min\_index = v;

return min\_index;

}

vector<vector<int>> MST(int parent[], int graph[V][V])

{

vector<vector<int>> v;

for (int i = 1; i < V; i++)

{

vector<int> p;

p.push\_back(parent[i]);

p.push\_back(i);

v.push\_back(p);

p.clear();

}

return v;

}

// getting the Minimum Spanning Tree from the given graph

// using Prim's Algorithm

vector<vector<int>> primMST(int graph[V][V])

{

int parent[V];

int key[V];

// to keep track of vertices already in MST

bool mstSet[V];

// initializing key value to INFINITE & false for all mstSet

for (int i = 0; i < V; i++)

key[i] = INT\_MAX, mstSet[i] = false;

// picking up the first vertex and assigning it to 0

key[0] = 0;

parent[0] = -1;

// The Loop

for (int count = 0; count < V - 1; count++)

{

// checking and updating values wrt minimum key

int u = minimum\_key(key, mstSet);

mstSet[u] = true;

for (int v = 0; v < V; v++)

if (graph[u][v] && mstSet[v] == false && graph[u][v] < key[v])

parent[v] = u, key[v] = graph[u][v];

}

vector<vector<int>> v;

v = MST(parent, graph);

return v;

}

// getting the preorder walk of the MST using DFS

void DFS(int\*\* edges\_list,int num\_nodes,int starting\_vertex,bool\* visited\_nodes)

{

// adding the node to final answer

final\_ans.push\_back(starting\_vertex);

// checking the visited status

visited\_nodes[starting\_vertex] = true;

// using a recursive call

for(int i=0;i<num\_nodes;i++)

{

if(i==starting\_vertex)

{

continue;

}

if(edges\_list[starting\_vertex][i]==1)

{

if(visited\_nodes[i])

{

continue;

}

DFS(edges\_list,num\_nodes,i,visited\_nodes);

}

}

}

int main()

{

// initial graph

int graph[V][V] = { { 0, 10, 18, 40, 20 },

{ 10, 0, 35, 15, 12 },

{ 18, 35, 0, 25, 25 },

{ 40, 15, 25, 0, 30 },

{ 20, 13, 25, 30, 0 } };

vector<vector<int>> v;

// getting the output as MST

v = primMST(graph);

// creating a dynamic matrix

int\*\* edges\_list = new int\*[V];

for(int i=0;i<V;i++)

{

edges\_list[i] = new int[V];

for(int j=0;j<V;j++)

{

edges\_list[i][j] = 0;

}

}

// setting up MST as adjacency matrix

for(int i=0;i<v.size();i++)

{

int first\_node = v[i][0];

int second\_node = v[i][1];

edges\_list[first\_node][second\_node] = 1;

edges\_list[second\_node][first\_node] = 1;

}

// a checker function for the DFS

bool\* visited\_nodes = new bool[V];

for(int i=0;i<V;i++)

{

bool visited\_node;

visited\_nodes[i] = false;

}

//performing DFS

DFS(edges\_list,V,0,visited\_nodes);

// adding the source node to the path

final\_ans.push\_back(final\_ans[0]);

// printing the path

cout<<"Optmial Path to travel: ";

for(int i=0;i<final\_ans.size();i++)

{

cout << final\_ans[i] << "-";

}

return 0;

}

\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*OUTPUT\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*

Optimal Path to travel: 0-1-3-4-2-0-

**Time Complexity:**

* The time complexity for obtaining MST from the given graph is O(V^2) where V is the number of nodes.
* The time complexity for obtaining the DFS of the given graph is O(V+E) where V is the number of nodes and E is the number of edges.
* Hence the overall time complexity is O(V^2).

**Space Complexity:**

* The worst case space complexity for the same is O(V^2), as we are constructing a vector<vector<int>> data structure to store the final MST.
* The space complexity for the DFS is O(V).
* Hence space complexity of this algorithm is O(V^2).

**Conclusion:**

**Hence studied travelling salesman problem and proposed a solution for it.**